

Optimality criterion for $\begin{cases} y \in Y \\ z \in Z \end{cases}$ is $\forall y \in Y \quad \nabla S_0(y)^T (y-y^*) \geq 0$

This is the first order optimality condition.

Amount moved
per unit cost change
min cost change as one moves from y^* to any y is positive if y^* is optimal

In our case the problem is:

$$\begin{cases} \min_z f_0(z) \\ \text{s.t. } \|z-x\|_2 \leq r \\ z \in C \end{cases}$$

$$\nabla S_0(z) = (z-x)$$

$$\nabla S_0(z^*) = (\pi_C(x)-x)$$

So, the optimality condition is $\forall z \in C \quad \nabla S_0(z^*)^T (z-z^*) = (\pi_C(x)-x)^T (z-\pi_C(x)) \geq 0$

Similarly when we are taking projection of y :

$$\forall z \in C \quad (\pi_C(y)-y)^T (z-\pi_C(y)) \geq 0$$

$$z = \pi_C(x)$$

$$(\pi_C(x)-x)^T (\pi_C(y)-\pi_C(x)) = (-1)^T (x-\pi_C(x))^T (\pi_C(x)-\pi_C(y)) \geq 0$$

$$(\pi_C(y)-y)^T (\pi_C(x)-\pi_C(y)) \geq 0$$

$$\begin{aligned} & (\pi_C(y)-y)^T (\pi_C(x)-\pi_C(y)) + (x-\pi_C(x))^T (\pi_C(x)-\pi_C(y)) \geq 0 \\ & = ((\pi_C(y)-y+x-\pi_C(x))^T (\pi_C(x)-\pi_C(y))) \\ & = (x-y)^T - (\pi_C(x)-\pi_C(y))^T (\pi_C(x)-\pi_C(y)) \\ & = (x-y)^T (\pi_C(x)-\pi_C(y)) - \|\pi_C(x)-\pi_C(y)\|_2^2 \geq 0 \\ & \rightarrow \|\pi_C(x)-\pi_C(y)\|_2^2 \leq (\pi_C(x)-\pi_C(y))^T (x-y) \leq \|\pi_C(x)-\pi_C(y)\|_2 \|x-y\|_2 \end{aligned}$$

By Cauchy-Schwarz
 $a^T b \leq \|a\| \|b\|$

$$\|\pi_C(x)-\pi_C(y)\|_2 \leq \|x-y\|_2$$

So $\pi_C(\cdot)$ is a Nonexpansive Operator

Similarly, overprojection operator $B_C = 2\pi_C - I$ on $C \subseteq \mathbb{R}^n$ is nonexpansive.

Caution: nonexpansive and contraction mapping are function, but monotone or strongly monotone operator can be nontrivial relations.
So monotone operator always splits out vector of same dimension as the input argument.

Monotone operators: def. monotone operators definitions and related

F relation: $R \subseteq \mathbb{R}^n \times \mathbb{R}^n$ monotone $\Leftrightarrow \forall (u,v) \in R, (x,y) \in R \quad (u-v)^T (x-y) \geq 0$

def. monotone operator

for $F(x) = x^*$ $(x^*-x^*)^T (x-y) = 0$

same x, y then $x^* > y^*$

Maximality: # of relations that are monotone operator is not finite

F maximal monotone $\Leftrightarrow (\exists \text{ monotone operator } F \subseteq F) \quad \nabla (u,v) \in F \Leftrightarrow \nabla (x,y) \in F$

def. maximal monotone

because fixing a relation is a set of $\{(u, F(u))\}$ structure

|| relation \rightarrow left just monotone

|| additionally \leftarrow right maximally monotone

Strong monotonicity:

F strongly monotone $\Leftrightarrow \exists m > 0 \quad \forall (u,v), (x,y) \in F \quad (u-v)^T (x-y) \geq m \|x-y\|_2^2$

$\forall x, y \in \text{dom } F \quad (F(x)-F(y))^T (x-y) \geq m \|x-y\|_2^2$ by Cauchy-Schwarz, $(F(x)-F(y))^T (x-y) \leq \|F(x)-F(y)\|_2 \|x-y\|_2 \leq L \|x-y\|_2^2$

|| $m < L$ ||

F strongly monotone Lipschitz with constant $L \geq \frac{1}{m}$

$\Rightarrow K = \frac{1}{m} \geq 1$

condition number of strongly monotone Lipschitz relation F.

Basic properties of monotone operator:

Sum and scalar multiple

F monotone, G monotone $\Rightarrow F+G$ is monotone [eq. sum of monotone is monotone]

F maximal monotone, G maximal monotone $\Rightarrow F+G$ maximal monotone [Additivity of maximal monotone operator with fine print]

F maximal monotone $\Rightarrow K F$ maximal monotone [Positive scalar multiplicativity of (maximal) monotone operator]

F strongly monotone with parameter m_1

G " " " " " " m_2

\Rightarrow

$(F+G)$ " " " " " " m_1+m_2

F strongly monotone $\Rightarrow K$ parameter m

F strongly monotone with parameter m_2

Inverse: which always exists

F maximal monotone $\Rightarrow F^{-1}$ maximal monotone [eq. inverse of monotone = monotone]

F strongly monotone with parameter $m_2 \Rightarrow F^{-1}$ function with Lipschitz constant $L = \frac{1}{m_2}$

What does that mean? It means that: inverse of a strongly monotone operator (which may be a nontrivial relation) is a function (as any relation with a Lipschitz constant is a function) which is freaking amazing.

Proof:

$\forall (x,u) \in F \quad \forall (z,v) \in F$

Cauchy-Schwarz

$\|u-v\|_2^2 \leq (u-v)^T (u-v) \leq \|u-v\|_2 \|u-v\|_2 \leq \|u-v\|_2 \|v-u\|_2$

definition of strong monotonicity

if we have:

Definition of monotone operator: let us start with scalar function

For that case monotone function is if input argument increases, the function value increases, and vice versa one way of quantifying that is:

$$(f(x)-f(y))(x-y) \geq 0 \Leftrightarrow \begin{cases} (x-y) > 0 \rightarrow (f(x)-f(y)) \geq 0 \\ (x-y) < 0 \rightarrow (f(x)-f(y)) \leq 0 \end{cases}$$

Note for imposing that as a definition we want that to be true for all (x,y) pairs, so we can give the definition:

$$\forall x \in \text{dom } f \Leftrightarrow \forall y \in \text{dom } f \quad (f(x)-f(y))(x-y) \geq 0$$

(+) direction of quantifier, (-) direction is interpretation of it. Note that the meaning is: $\forall x \in \text{dom } f \quad (f(x)-f(y))(x-y) \geq 0 \Leftrightarrow \forall x \in \text{dom } f$

Now let's extend this for vector valued case, as vectors are partial ordered we cannot just start with $x > y \Rightarrow f(x) > f(y)$, however look at the scalar case $\forall (u,v) \in (F(x)-F(y))(x-y) \geq 0$ and trying to extend it point wise for a vector case: $\forall x, y \in \mathbb{R}^n \quad (F(x)-F(y))^T (x-y) = (F(x)-F(y))^T (x-y) \geq 0$

What might be an intuition of this? Well it means that in some element of the vector the monotone operator is results in a such scalar monotonic behavior to balance any non-monotonic behavior in other elements. For a 2D case: $(f(x)-f(y))^T (x-y) = (f(x)-f(y))^T (x-y) \geq 0$

So if the first term gets negative for some x,y , then the other terms will counter it by getting strongly positive. Note that this means:

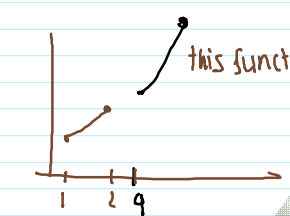
$$\forall (x,y) \in F \Leftrightarrow \forall (z,y) \in F \quad (u-v)^T (x-y) \geq 0$$

-this is monotone, if we modify the notation as follows:

$$\forall (x,u) \in F \Leftrightarrow \forall (z,v) \in F \quad (u-v)^T (x-y) \geq 0$$

maximal monotone

For example



this function is: $f(x) = x^2, \text{ dom } f = [1, 2] \cup [4, \infty)$, so if a function is not maximal monotone then

now set $x = 3 \notin \text{dom } f$

$\forall y \in \text{dom } f = [1, 2] \cup [4, \infty)$

$$(f(x)-f(y))(x-y) = (9-y^2)(3-y) = (3+y)(3-y)(3-y) = (3+y)(3-y)^2 \geq 0$$

$\therefore \exists x \in \mathbb{R} \quad \forall y \in \text{dom } f \quad (f(x)-f(y))(x-y) \geq 0 \Leftrightarrow x \notin \text{dom } f$

\therefore So, f is not maximal monotone

But it can be shown that, $\forall x \in \text{dom } f \rightarrow \forall y \in \text{dom } f \quad (f(x)-f(y))(x-y) = (x^2-y^2)(x-y) = (x-y)(x+y)(x-y) \geq 0$ i.e., f is monotone trivially

\therefore So, in summary we have a function $f(x)$ which is monotone but not maximal monotone

$\forall x \in \text{int } C, N_C(x) = \{0\}$
 $\forall x \in \partial C, N_C(x) = \text{nontrivial}$ # The intuition behind this is that remember at the boundary of any function, the subdifferential might not exist. One strange function in this regard is the function of \mathbb{R}^2 set.
 Why $N_C(x)$ matters? Because this is the subdifferential mapping of the convex indicator function $\mathbb{1}_C(\cdot) = \begin{cases} 0, & \exists \in C \\ \infty, & \exists \notin C \end{cases}$, i.e., $N_C = \partial \mathbb{1}_C$.
 Makes sense because remember, for any convex function the subdifferential always exists in the interior of f . At the boundary the subdifferential may not exist, i.e., $\forall x \in \partial(C) \partial \mathbb{1}_C(x)$ might not exist, as a result $N_C(x \in \partial C) = \partial \mathbb{1}_C(x \in \partial C) = \text{nontrivial}$.
 But remember for all the interior points, $N_C(x) = \partial \mathbb{1}_C(x) = \{0\}$.

[eq: subdifferential is monotone mapping]
 Proof: $\mathbb{1}_C(x) = \begin{cases} 0, & x \in C \\ \infty, & x \notin C \end{cases}$. Want to determine, $\partial \mathbb{1}_C(x) = N_C(x)$.
 By definition of subdifferential, $g \in \partial \mathbb{1}_C(x) \iff \forall y \in \mathbb{R}^n, \mathbb{1}_C(y) \geq \mathbb{1}_C(x) + g^T(y-x) \iff \forall y \in \mathbb{R}^n, \mathbb{1}_C(y) - \mathbb{1}_C(x) \geq g^T(y-x)$.
 • Case 1: $x \in C$, then $\forall y \in C, 0 \geq g^T(y-x) \iff \forall y \in C, g^T(y-x) \leq 0$.
 $\therefore x \in C \implies \partial \mathbb{1}_C(x) = \{g \mid \forall y \in C, g^T(y-x) \leq 0\}$.
 • Case 2: $x \notin C \implies 0 \geq \infty + g^T(y-x) \implies$ no finite g can exist $\therefore \partial \mathbb{1}_C(x) = \emptyset$.
 $\therefore \partial \mathbb{1}_C(x) = N_C(x) = \begin{cases} \emptyset, & x \notin C \\ \{g \mid \forall y \in C, g^T(y-x) \leq 0\}, & x \in \text{int}(C) \\ \text{nontrivial}, & x \in \partial C \end{cases}$

Not the most rigorous proof...

• Saddle subdifferential:
 $f: \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R} \cup \{+\infty, -\infty\}$, assume f is convex in θ , concave in x .

$f(x, y) = \{ \text{saddle subdifferential relation} \}$
 $\begin{bmatrix} \partial_x f(x, y) \\ \partial_y (-f(x, y)) \end{bmatrix} \iff \begin{cases} \partial_x f(x, y) \neq \emptyset, \partial_y f(x, y) \neq \emptyset \\ -f(x, y) : \text{convex in } \theta \end{cases}$

Why this is called saddle subdifferential?
 Because $(x, y) \in F \iff (x, y) \in \text{argmin}_{\theta} f(x, \theta) \iff \exists \lambda \in \mathbb{R}^m, \begin{bmatrix} \partial_x f(x, y) \\ \partial_y (-f(x, y)) \end{bmatrix} \ni \begin{cases} 0 \in \partial_x f(x, y) \iff x \in \text{argmin}_{x'} f(x', y) \iff \forall (x', y) \in S(x, y) \\ 0 \in \partial_y (-f(x, y)) \iff y \in \text{argmax}_{y'} (-f(x, y')) \iff \forall (x, y') \in S(x, y) \end{cases}$
 $\implies \forall (x', y') \in S(x, y), f(x, y) \leq f(x', y) \leq f(x, y')$
 \therefore Set of F is the saddle points of f .

Saddle subdifferential relation for CP function $f(x, y)$ is maximal

* KKT operator: # Manuscript version

$\begin{pmatrix} \nabla f(x, y) \\ \nabla g_1(x, y) \\ \vdots \\ \nabla g_m(x, y) \\ \nabla h_1(x, y) \\ \vdots \\ \nabla h_l(x, y) \end{pmatrix} \begin{matrix} \lambda \\ \mu \\ \vdots \\ \nu \end{matrix}$
 $L(x, \lambda, \mu, \nu) = \begin{cases} f(x, y) + \sum_{i=1}^m \lambda_i g_i(x, y) + \sum_{j=1}^l \nu_j h_j(x, y), & \text{if } \lambda \geq 0 \\ -\infty, & \text{otherwise} \end{cases}$
 $T(x, \lambda, \mu, \nu) = \text{KKT operator} = \begin{bmatrix} \partial_x L(x, \lambda, \mu, \nu) \\ -\partial_y L(x, \lambda, \mu, \nu) \\ \partial_x g_1(x, y) \\ \vdots \\ \partial_x g_m(x, y) \\ \partial_x h_1(x, y) \\ \vdots \\ \partial_x h_l(x, y) \end{bmatrix}$, so $(x, \lambda, \mu, \nu)^T \in T \iff \begin{cases} -\partial_x L(x, \lambda, \mu, \nu) = [-s; \text{ineq}] + \partial_x g_i(x, y) = 0 \iff \text{primal infeasibility} \\ -\partial_y L(x, \lambda, \mu, \nu) = [-s; \text{ineq}] = 0 \iff \text{dual feasibility condition} \end{cases}$

$T(x, \lambda, \mu, \nu)$ Special case of saddle subdifferential = [maximize]

$T(x^*, \lambda^*, \mu^*, \nu^*) \ni T(x^*, \lambda^*, \mu^*, \nu^*)$ solves the minimization problem

* KKT operator (Simpler version)

$\forall f(x)$
 $\exists Ax = b$
 $L(x, y) = f(x) + y^T(Ax - b)$
 KKT operator: $F(x, y) = \begin{bmatrix} \partial_x L(x, y) \\ -\partial_y L(x, y) \end{bmatrix} \begin{matrix} \lambda \\ \mu \end{matrix}$ # gives vanishing gradient at Lagrangian condition (one of the KKT conditions)
 # gives LHS of primal feasibility (another KKT condition)

Note that this is an equality constrained optimization problem, so it will have only these two KKT conditions with RHS zero and the LHS 1 given by the rows of the KKT operator. So if the optimal (x^*, y^*) is inputted through the KKT operator then we will have zero.

$F(x, y) = \begin{bmatrix} \partial_x L(x, y) \\ -\partial_y L(x, y) \end{bmatrix}$

so if $F(x^*, y^*) = 0 \iff (x^*, y^*)$: optimal primal dual pair

- $\iff (x^*, y^*, 0) \in F$
- $\iff (0, (x^*, y^*)) \in F^{-1}$
- $\iff (x^*, y^*) \in F^{-1}(0)$

So optimal primal dual pair will belong to the output set of the inverse KKT operator with 0 fed into it!

Multiplier to residual mapping is very important as it has a connection with ADMM

* Multiplier to residual mapping: def: multiplier to residual mapping

$\begin{pmatrix} \nabla f(x) \in \partial f(x) \\ Ax = b \end{pmatrix}$

$L(x, y) = f(x) + y^T(Ax - b)$ Inner technicals: $F(x) = b - A \text{argmin}_x L(x, y) \therefore F(x) = b - A \text{argmin}_x [f(x) + y^T(Ax - b)] = b - A \text{argmin}_x [f(x) + y^T Ax - y^T b]$

Assume $F(x) = b - A x^*(y)$ # $x^*(y) = \text{argmin}_x L(x, y)$, because $f(x)$: strongly convex, $y^T(Ax - b) = Ay^T x - y^T b$: affine in $x \implies f(x) + y^T(Ax - b)$: strongly convex, so will have unique minimizer in x . # Note: convex + strongly convex = convex.

This F is called multiplier to residual mapping because it takes the Lagrangian multiplier and outputs the residual that associated with the sub-optimality of x .
 $\partial_x L(x, y) = 0 \iff \partial_x f(x) + A^T y = 0 \iff \partial_x f(x) = -A^T y$
 $\partial_x [f(x) + y^T(Ax - b)] = \partial_x f(x) + A^T y = 0$
 $\iff \exists x \in \text{argmin}_x [f(x) + y^T(Ax - b)]$
 $\iff -A^T y \in \partial_x f(x) \iff (x, -A^T y) \in \partial_x f$
 $\iff x^*(y) = (x, -A^T y) \in (\partial_x f)^{-1}(0)$

Alternative definition to multiplier is # residual mapping operator

$F(y) = b - A (\partial_x f)^{-1}(A^T y)$ this is a monotone operator, so multiplier to residual mapping is a monotone operator

proof: $(\partial_x f)^{-1}$ is monotone [eq: subdifferential is monotone mapping]
 $\implies (\partial_x f)^{-1}$ is monotone [eq: inverse of monotone = monotone]
 $\implies \forall y_1, y_2, \exists x_1, x_2, f_1^{-1}(y_1) \preceq f_1^{-1}(y_2) \iff$ [eq: composition of monotone = monotone] # F : monotone $\rightarrow A^T F(A \cdot)$: monotone
 \downarrow
 $\text{III} = -A^T$
 $-A (\partial_x f)^{-1}(A^T \cdot) \preceq -A (\partial_x f)^{-1}(A^T \cdot)$
 $\implies (b - A (\partial_x f)^{-1}(A^T y_1)) \preceq (b - A (\partial_x f)^{-1}(A^T y_2))$

Addition of a constant vector to each element of a relation set will not change maximality of a relation; equivalent logic. $\mathbb{1} \oplus \mathbb{B}$ is a monotone operator
 $\mathbb{B} \oplus (-A)(\mathbb{1} \oplus \mathbb{1})^{-1}(-A \oplus \mathbb{B})$ is a monotone operator
 $-A(\mathbb{1} \oplus \mathbb{1})^{-1}(-A \oplus \mathbb{B})$ is a monotone operator
(set form of monotone is important)

Now we show:

$$\begin{aligned}
 * F(y) &= b - A(\partial f)^{-1}(-A^T y) \\
 &= \partial y [b^T y + f^*(-A^T y)]
 \end{aligned}$$

Proof:

$$\partial y [b^T y + f^*(-A^T y)] = \partial y \underbrace{[f^*(-A^T y)]}_{-A \underbrace{[\partial f^*(y)]}_{y = -A^T y}}$$

/* chain rule: $\partial_x f(Ax) = A^T [\partial_x f(x)]_{x=Ax}$ */

Now $\partial y f^*(y) = \partial y \left[\sup_x -f(x) + y^T x \right]$

$$\begin{aligned}
 &= \partial y [-f(x^*) + y^T x^*] = \partial y \underbrace{[x^*]}_{x^*} = \text{argmax}_x -f(x) + y^T x \quad \left\{ \begin{array}{l} \text{+ recall, if } f = \sup_{K \in A} (f_K) \\ \text{if } f = \sup_{K \in A} f_K = f_K \end{array} \right\} \\
 &= \text{argmin}_x f(x) - y^T x \quad \left\{ \begin{array}{l} \text{find } \eta \in \partial f \Rightarrow \eta \in \partial f^* \text{ !} \\ \text{find } \eta \in \partial f \Rightarrow \eta \in \partial f^* \text{ !} \end{array} \right\}
 \end{aligned}$$

$x^* = \text{argmax}_x [-f(x) + y^T x] = -f(x^*) + y^T x^* = \text{argmax}_x -f(x) + y^T x$
 $= \text{argmin}_x f(x) - y^T x$
 $\Leftrightarrow [\partial_x (f(x) - y^T x)]_{x=x^*} \ni 0$
 $\Leftrightarrow \partial f(x^*) - y \ni 0 \Leftrightarrow \partial f(x^*) \ni y$
 $\Leftrightarrow (x^*, y) \in \partial f \Leftrightarrow (y, x^*) \in (\partial f)^{-1} \Leftrightarrow x^* = (\partial f^{-1})(y) \quad \exists!$
This is a function (inverse of subdifferential is a function)

$\therefore \partial y f^*(y) = (\partial f^{-1})(y)$

$$\begin{aligned}
 \therefore \partial y [b^T y + f^*(-A^T y)] &= b - \partial y [f^*(-A^T y)] = b - [(\partial f^{-1})(y)]_{y = -A^T y} \\
 &= b - A(\partial f^{-1})(-A^T y) = b - A \underbrace{[\partial f^{-1}(y)]}_{y = -A^T y} = b - A(\partial f^{-1})(-A^T y) \quad \blacksquare
 \end{aligned}$$